

The Benefits of Solid State in Enterprise Storage Systems

David Dale, NetApp

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Abstract



Solid State in Enterprise Storage Systems

- * Targeted primarily at an IT audience, this session presents a brief overview of the solid state technologies which are being integrated into Enterprise Storage Systems today, including technologies, benefits, and price/performance.
- It then goes on to describe where they fit into typical Enterprise Storage architectures today, with descriptions of specific use cases.
- Finally the presentation speculates briefly on what the future will bring.

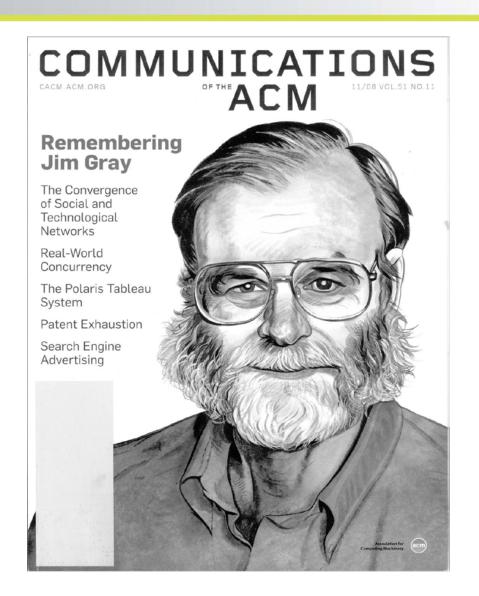
Agenda



- Why flash in the datacenter? Why now?
- Memory, cache and storage
- Application opportunities
- Flash in enterprise storage today
 - SSD storage tier
 - Network cache
 - Storage controller-based cache
- What's next
- Conclusion

Remembering Jim Gray





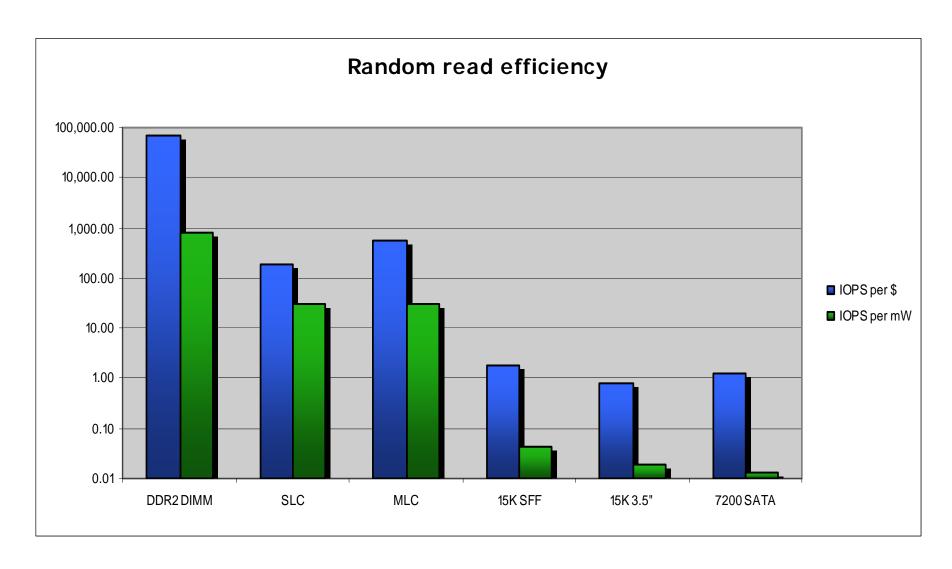
Database and systems design pioneer, and cocreator of the Five Minute Rule (1987)

"Flash is a better disk ..., and disk is a better tape" ~2006

Lost at sea January 2007

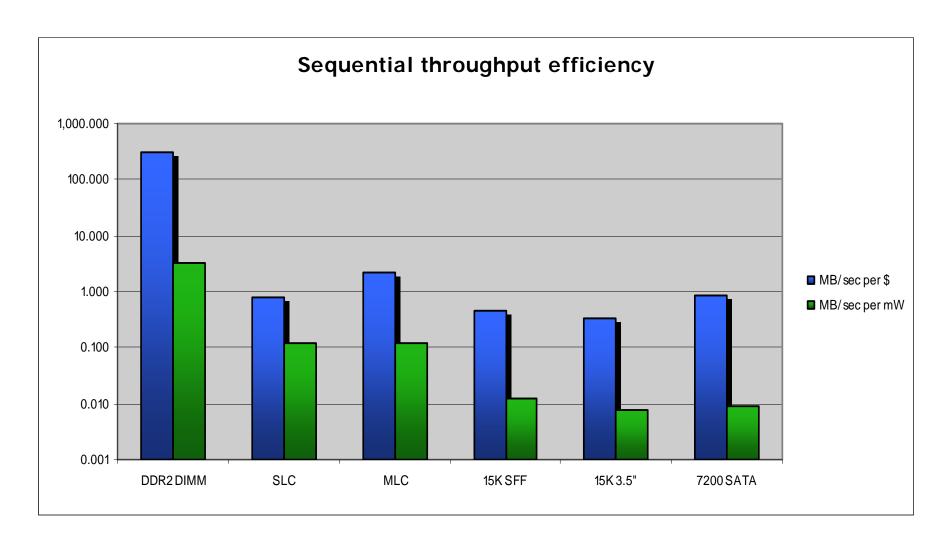
Why Flash? IOPS efficiency vs. HDD





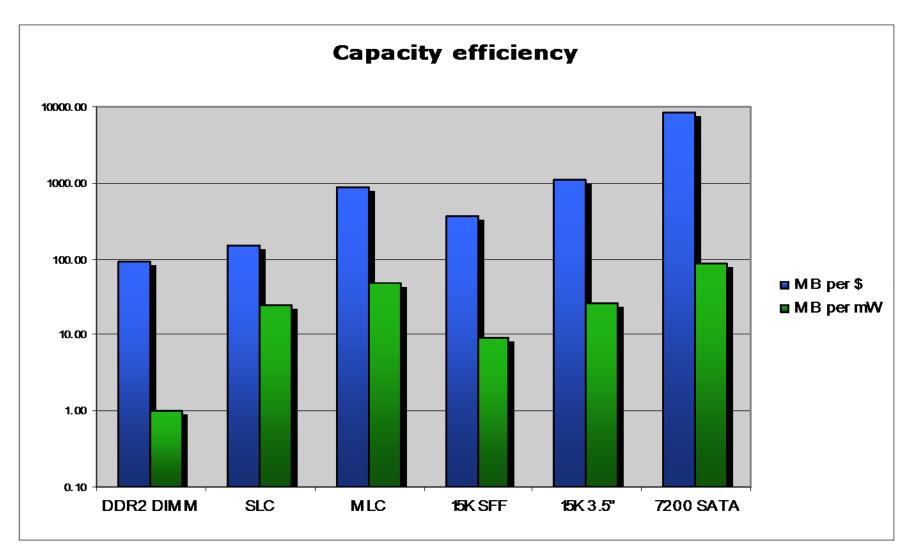
Why Flash? Bandwidth/Watt vs. HDD





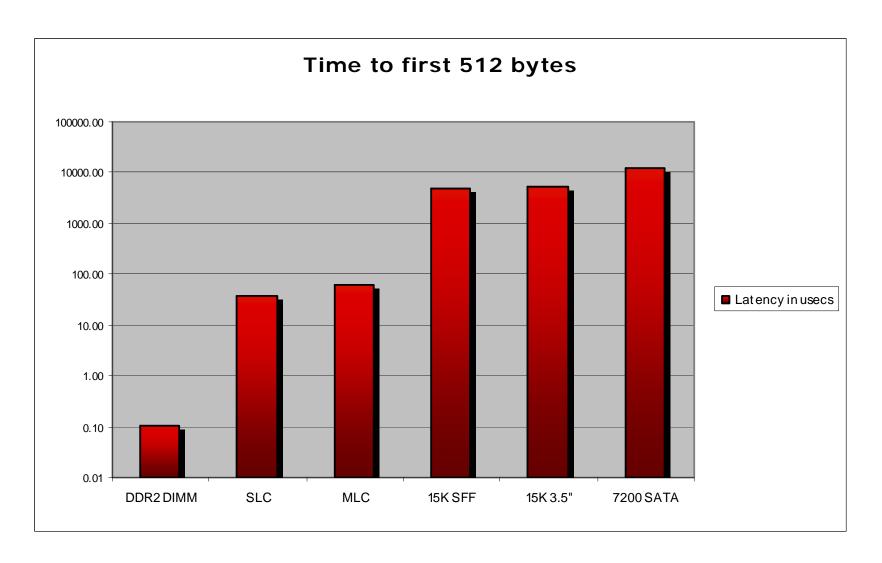
Why Flash? Capacity/Watt vs. DRAM





Why Flash? Latency vs. HDD





Why Flash in the Datacenter Now?



Why flash?

- Capacity efficiency versus DRAM
 - > ~10x better \$ per GB
 - > ~30x better power per GB
- IOPS efficiency versus HDDs
 - > ~100x better \$ per IOPS
 - > ~1000x better power per IOPS

Why now?

- Period of rapid density advancements led to HDD-like bit density at lower \$/GB than DRAM
- Innovations in SSD and tiering technology

Five-Minute Rule, 1987



Assuming that the cost of a cache is dominated by its capacity, and the cost of a backing store is dominated by its access cost (cost per IOPS), then the breakeven interval for keeping a page of data in cache is given by:

Break-Even-Interval =

Backing-Store-Cost-Per-IOPS /

Cache-Cost-Per-Page

♦ 1987: Disk \$2,000 / IOPS; RAM \$5 / KB →
I KB breakeven = 400 seconds ~= 5 minutes

Five-Minute Rule, 2008



- ◆Disk \$1 / IOPS (2,000x reduction)
- ◆DRAM \$50 / GB (100,000x reduction) \$0.05 / MB, \$0.0025 / 50 KB
- → 50 KB breakeven ~= 5 minutes
- → 4 KB breakeven ~= 1 hour
- → I KB breakeven ~= 5 hours as Gray predicted
- 100,000x / 2,000x = 50-fold increase in size of "page" to cache for breakeven at 5 minutes

Five-Minute Rule, 2010: Flash & HDD



- →HDD \$1 / IOPS (2,000x reduction)
- SLC flash ~\$10 / GB (packaged)
- ♦ MLC flash ~\$4 / GB (packaged)
- → SLC 250 KB breakeven ~= 5 minutes SLC 4 KB breakeven ~= 5 hours
- → MLC 625 KB breakeven ~= 5 minutes MLC 4 KB breakeven ~= 13 hours
- 100,000x / 2,000x = 50-fold increase in size of "page" to cache for breakeven at 5 minutes

Five-Minute Rule, 2010: DRAM & Flash



- ◆SLC flash ~\$0.05 / IOPS (4 KB, enterprise SSD)
- ♦ MLC flash ~\$0.02 / IOPS (4 KB, enterprise SSD)
- ◆DRAM \$20 / GB (enterprise DIMMs)
- → SLC 6 KB breakeven ~= 5 minutes
- → MLC 2 KB breakeven ~= 5 minutes

Need to consider however that cost/capacity of flash (at least SLC) is a large fraction of that of DRAM

Implications



- → Flash makes it cost-effective to keep more small random data in silicon-based cache versus DRAM:
 - ~5+ hour working set versus ~I hour
- Flash allows small random data working set in DRAM to be reduced, allowing cost, power, space efficiency:
 minute working set versus ~I hour
- Assuming appropriate locality of reference, transfer sizes between HDD and flash tiers should increase to preserve expensive HDD IOPS
- ❖Flash tier likely to alter checkpoint processing intervals (shorter), metadata organization (e.g. optimal page size)

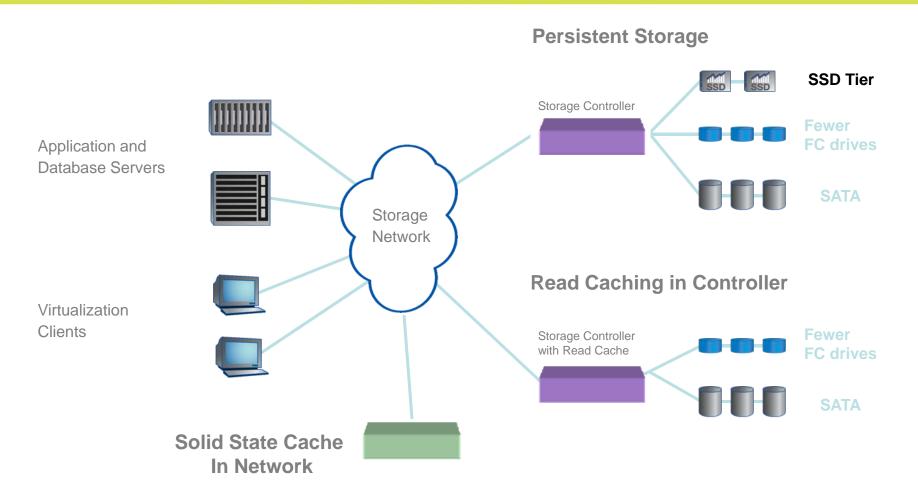
Application opportunities



- > Intense random reads, e.g. OLTP, metadata
- Sequential read after random write
 - Log-oriented writes convert this to random read after sequential write (e.g. FTL)
- Low read latency (~100x better than HDD)
 - Facilitates DRAM extension by allowing high read throughput with limited read concurrency
 - Paging datacenter apps can be practical again
 - Memory capacity to consolidate more servers with underutilized CPU
- Enabling memory-resident datasets, e.g.
 - OLTP
 - Data warehouses (viz TPC-H results)
 - Large metadata

Storage Networking with Flash





Available Solutions: Pro and Con



| Technology | Pros | Cons |
|-----------------------|---|--|
| Solid State Drives | Response times consistently fast for reads Low cost per IOP Administrator has direct control over data stored in SSD tier | High cost per gigabyte Requires software tools and administration to move hot data into and out of SSD tier Limited apps today |
| Controller Read Cache | Hot data automatically flows into read cache—no administration required Deployment is relatively non-disruptive Viable for common enterprise applications | Cache must be populated before it becomes effective |
| Network Cache | Hot data automatically flows into the caching tier Deployment is relatively non-disruptive Scalable solution for high performance applications | Cache must be populated before it becomes effective May be limited by protocol choice available from vendor |

SSD Is Much Faster than HDD





SSD tier



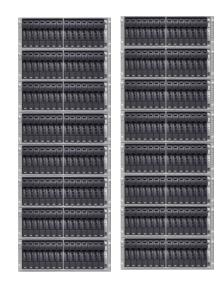
HDD tier

| | SDD | HDD | Comment |
|------------------|-------------|-------------|-------------------------------------|
| Capacity | Equal (2TB) | Equal (2TB) | Hold capacity constant |
| IOPs | ~50,000 | ~3,600 | Order of magnitude faster for flash |
| Latency | Better | Worse | Order of magnitude lower for flash |
| Carbon Footprint | Better | Worse | Same per TB, better per IOPs |

For a 50,000 IOP System, SSD Has Much Betterion In Latency and Footprint



SSD tier



HDD Tier

| | SSD Tier | HDD Tier | Comment |
|------------|--------------|----------------|------------------------------------|
| IOPs | 50,000 | 50,000 | Hold IOPs constant |
| Capacity | Worse (2TB) | Better (27 TB) | Significantly more for disk |
| Latency | Better (1ms) | Worse (10ms) | Order of magnitude lower for flash |
| Rack Space | 12U | 54U | Significantly more for disk |

SSD Tier



Advantages:

- Fast random I/O for small blocks
- Low read and write latency time
- Low power consumption
- Low noise
- Better mechanical reliability

Disadvantages:

- Very high price, typically 10-30 X comparable FC drives
- Limited capacities
- Slow random write speeds, e.g. erase of blocks
- Slow sequential write throughput

SSD-based Solutions



Database acceleration solution

- Entire DB on SSD tier
- Or hot files on SSD and rest of DB on standard disk
 - > Redo logs, indexes, temp space

Large scale virtual machine environments

Solves "boot storm" problem for large numbers of virtual machines

Entire data set on SSD tier

- Multiple apps: Virtual Servers, VDI, and so on
- Any app where entire data set can fit into memory of SSD array
- Works well in NAS environments

Network cache solutions

- All files on HDD in shared storage array
- Accelerated by SSD-based network cache
- Self—tuning write-through cache
- Applications include
 - > Rendering, seismic, financial modeling, ASIC design

Controller-based Flash Cache



- Functions as an intelligent read cache for data and metadata
- Automatically places active data where access can be fast
- Provides more I/O throughput without adding high-performance disk drives to a disk-bound storage system
- Effective for file services, OLTP databases, messaging, and virtual infrastructure

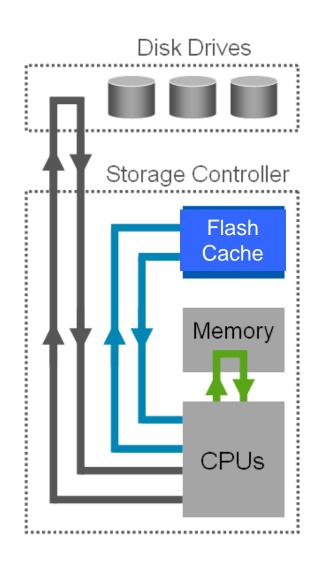
Deciding Between SSD and Cache

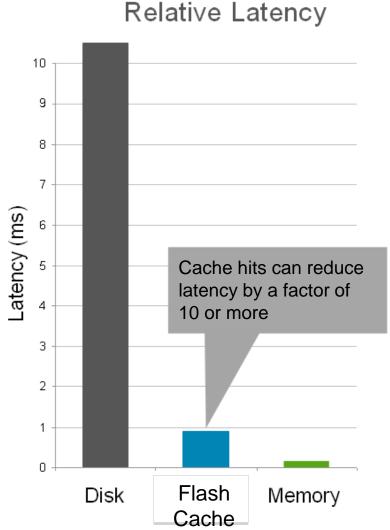


| SSD Persistent Storage | Controller-Based Read Caching |
|--|---|
| Good Fit When | Good Fit When |
| Random I/O intensive workload | Random read intensive workload |
| Every read must be fast | Improving average response time |
| Active data is known and fits | is adequate |
| into the SSD tier | Active data is unpredictable |
| Active data is known, is dynamic, | or unknown |
| and ongoing administration is OK | Administration-free approach is |
| Upside of write acceleration desired | desired |
| | Start small and scale up |

Reduce Latency with Flash Cache

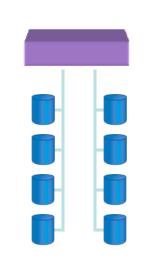






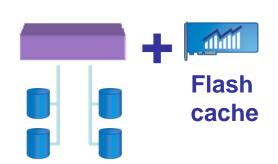
Configuring Storage with Flash Cache





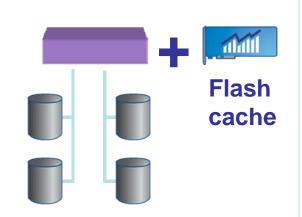
Configure with **FC** Disks Only

- Additional disk drives provide IOPs
- Inefficient use of storage capacity, power, and space



Configure with FC Disks And Flash Cache

- Disks provide capacity/IOPs
- Cache delivers more IOPs and speeds response times
- Achieve cost savings for storage, power, and space

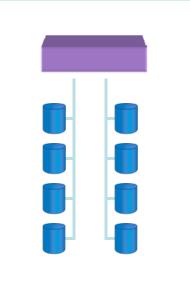


Configure with **SATA**Disks and **Flash Cache**

- More storage capacity
- Cache provides IOPs boost for SATA disk drives
- Achieve cost savings for storage, power, and space

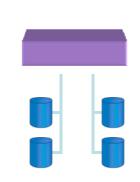
Use case: Scale Performance of Disk-bound Systems







- Use more disks to provide more IOPs
- May waste storage capacity
- Consumes more power and space





- Performance is disk-bound
- Have enough storage capacity
- Random read intensive workload





Add Flash Cache

- Use cache to provide more IOPs
- Improves response times
- Uses storage efficiently
- Achieves cost savings for storage, power, and space



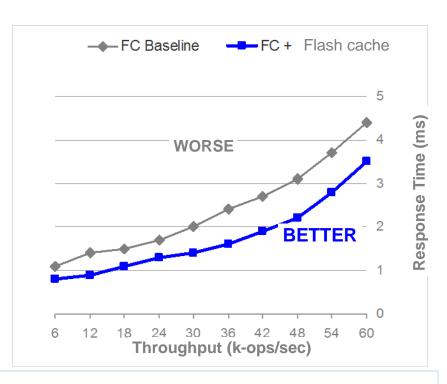
FC HDD plus Flash Cache Example



Benchmarked Configurations

224 FC drives 64TB 75% Fewer Spindles 56 FC drives 16TB FC Baseline Configuration FC + Flash Cache Configuration

SPECsfs2008 Performance



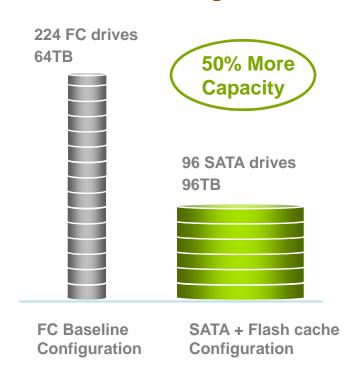
- Purchase price is 50% lower for FC + Flash cache compared to Fibre Channel baseline
- FC + Flash cache yields 67% power savings and 67% space savings

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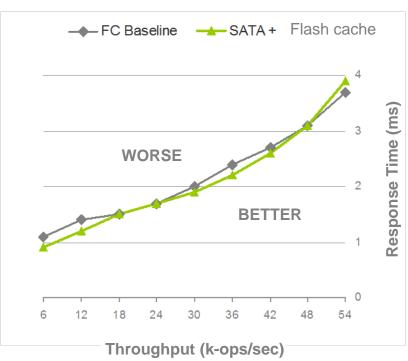
SATA HDD plus Flash Cache Example



Benchmarked Configurations



SPECsfs2008 Performance

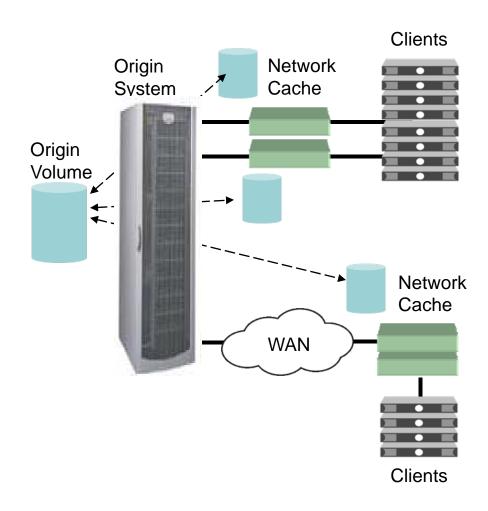


- Purchase price is 39% lower for SATA + Flash cache compared to Fibre Channel baseline
- SATA + Flash cache yields 66% power savings and 59% space savings

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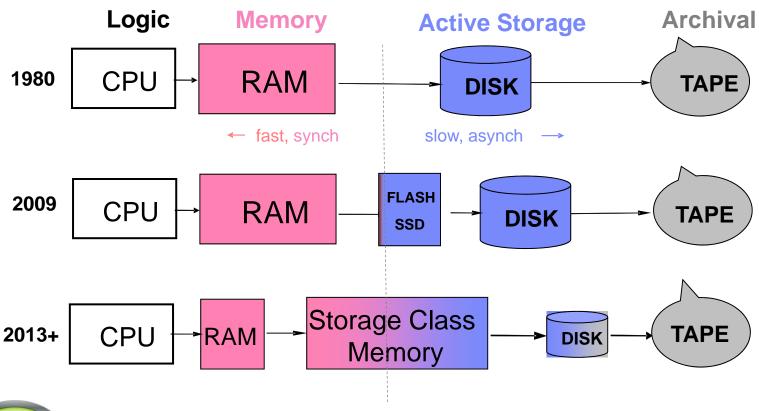
Network Cache Topology





System Evolution







Check out SNIA Tutorial:

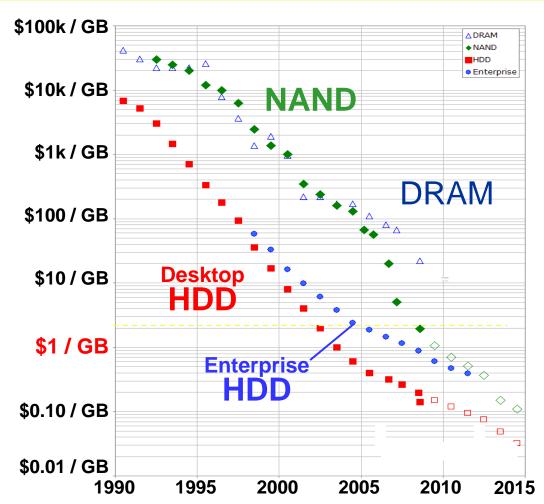
"The Future of Solid State Storage"

Cost Structure of Silicon-based Technology



Cost determined by

- cost per wafer
- # of dies/wafer
 - memory area
 per die [sq. μm]
 - memory density
 [bits per 4F²]
 - patterning density
 [sq. μm per 4F²]





Check out SNIA Tutorial: "The Future of Solid State Storage"

Chart courtesy of Dr. Chung Lam, IBM Research updated version of plot from 2008 IBM Journal R&D article

Summary



- Over the next 5 years solid state technologies will have a profound impact on enterprise storage
- It's not just about replacing mechanical media with solid state media
- The architectural balance of memory, cache and persistent storage will change
- Today's solid state implementations in enterprise storage demonstrate these changes
- It's only the beginning...

Q&A / Feedback



Please send any questions or comments on this presentation to SNIA: tracksolidstate@snia.org

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- SNIA Education Committee

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